

# Decomposition Applications

- rate control
- single commodity network flow

## Rate control setup

- $n$  flows, with fixed routes, in a network with  $m$  links
- variable  $f_j \geq 0$  denotes the rate of flow  $j$
- flow utility is  $U_j : \mathbf{R} \rightarrow \mathbf{R}$ , strictly concave, increasing
- traffic  $t_i$  on link  $i$  is sum of flows passing through it
- $t = Rf$ , where  $R$  is the routing matrix

$$R_{ij} = \begin{cases} 1 & \text{flow } j \text{ passes over link } i \\ 0 & \text{otherwise} \end{cases}$$

- link capacity constraint:  $t \preceq c$

# Rate control problem

$$\begin{array}{ll} \text{maximize} & U(f) = \sum_{j=1}^n U_j(f_j) \\ \text{subject to} & Rf \preceq c \end{array}$$

- convex problem
- dual decomposition gives decentralized method

## Rate control Lagrangian

Lagrangian (for minimizing  $-U$ ) is

$$\begin{aligned} L(f, \lambda) &= -U(f) + \lambda^T (Rf - c) \\ &= -\lambda^T c + \sum_{j=1}^n (-U_j(f_j) + (r_j^T \lambda) f_j) \end{aligned}$$

- $\lambda_i$  is price (per unit flow) for using link  $i$
- $r_j^T \lambda$  is the sum of prices along route  $j$

## Rate control dual

dual function is

$$\begin{aligned} g(\lambda) &= -\lambda^T c + \sum_{j=1}^n \inf_{f_j} (-U_j(f_j) + (r_j^T \lambda) f_j) \\ &= -\lambda^T c - \sum_{j=1}^n (-U_j)^*(-r_j^T \lambda), \end{aligned}$$

dual rate control problem:

$$\begin{aligned} &\text{maximize} && -\lambda^T c - \sum_{j=1}^n (-U_j)^*(-r_j^T \lambda) \\ &\text{subject to} && \lambda \succeq 0 \end{aligned}$$

subgradient of negative dual:

$$R\bar{f} - c \in \partial(-g)(\lambda)$$

where  $\bar{f}_j = \operatorname{argmax} (U_j(f_j) - (r_j^T \lambda) f_j)$

## Dual decomposition rate control algorithm

**given** initial link price vector  $\lambda \succeq 0$  (e.g.,  $\lambda = \mathbf{1}$ ).

**repeat**

1. Sum link prices along each route.  
Calculate  $\Lambda_j = r_j^T \lambda$ .
2. Optimize flows (separately) using flow prices.  
 $f_j := \operatorname{argmax} (U_j(f_j) - \Lambda_j f_j)$ .
3. Calculate link capacity margins.  
 $s := c - Rf$ .
4. Update link prices.  
 $\lambda := (\lambda - \alpha_k s)_+$ .

## Dual decomposition rate control algorithm

- decentralized:
  - links only need to know the flows that pass through them
  - flows only need to know prices on links they pass through
- prices converge to optimal; so do flows (since  $U$  is strictly concave)
- iterates can be (and often are) infeasible, *i.e.*,  $Rf \not\leq c$   
(but we do have  $Rf \leq c$  in the limit)
- have upper bound  $-g(\lambda)$  on optimal utility  $U^*$



## Generating feasible flows

- define  $\eta_i = t_i/c_i = (Rf)_i/c_i$ 
  - $\eta_i < 1$  means link  $i$  is under capacity
  - $\eta_i > 1$  means link  $i$  is over capacity
- define  $f^{\text{feas}}$  as

$$f_j^{\text{feas}} = \frac{f_j}{\max\{\eta_i \mid \text{flow } j \text{ passes over link } i\}}$$

- $f^{\text{feas}}$  will be feasible, even if  $f$  is not
- finding  $f^{\text{feas}}$  is also decentralized  
(in fact this is a step in primal decomposition)

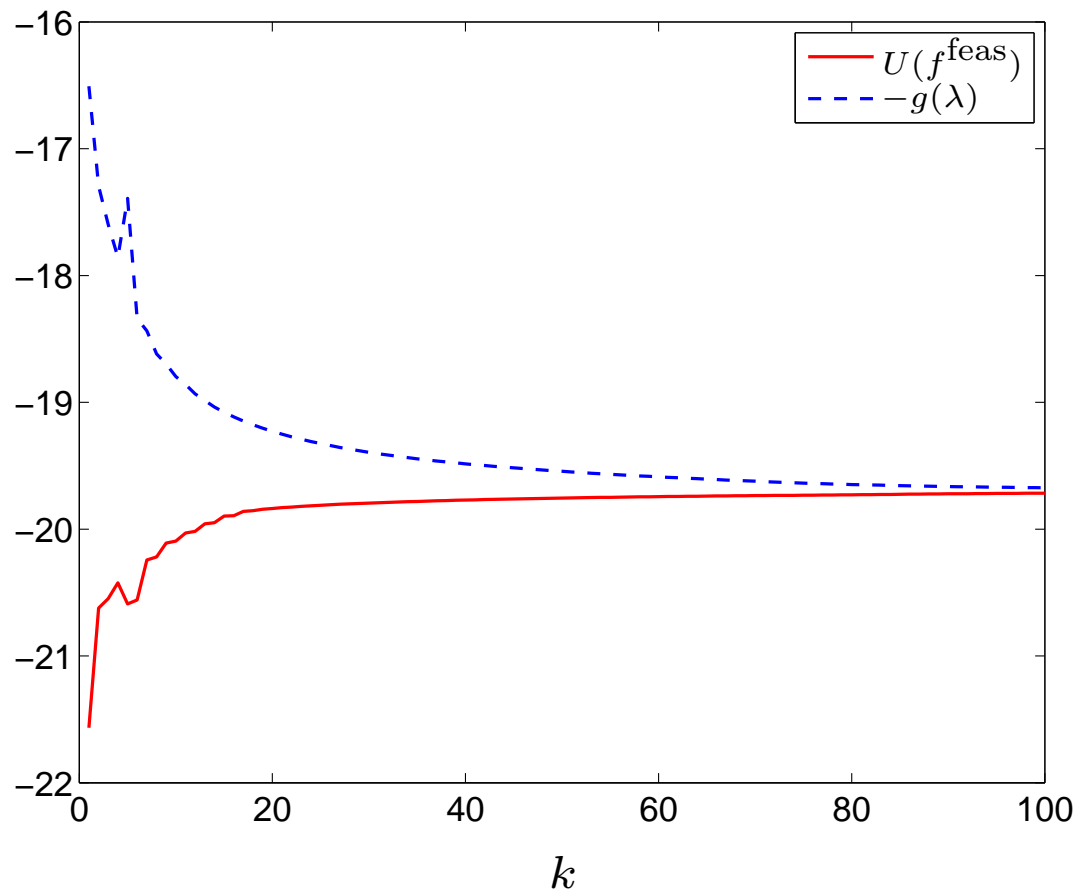
## Example

- $n = 10$  flows,  $m = 12$  links; 3 or 4 links per flow
- link capacities chosen randomly, uniform on  $[0.1, 1]$
- $U_j(f_j) = \log f_j$  (can be argued to give proportionally fair flows)
- optimal flow as a function of price:

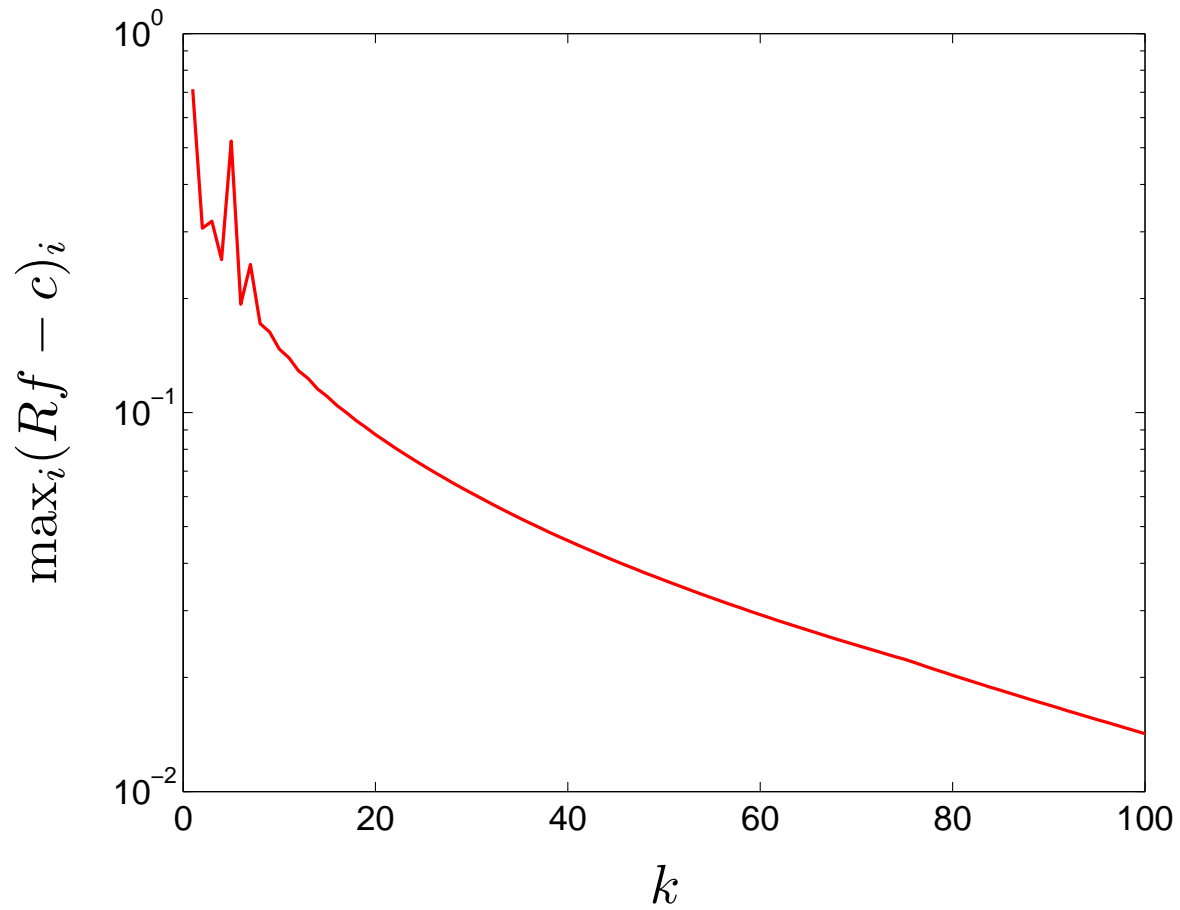
$$\bar{f}_j = \operatorname{argmax}(U_j(f_j) - \Lambda_j f_j) = 1/\Lambda_j$$

- initial prices:  $\lambda = \mathbf{1}$
- constant stepsize  $\alpha_k = 3$

# Convergence of primal and dual objectives



# Maximum capacity violation



## Single commodity network flow setup

- connected, directed graph with  $n$  links,  $p$  nodes
- variable  $x_j$  denotes flow (traffic) on arc  $j$
- given external source (or sink) flow  $s_i$  at node  $i$ ,  $\mathbf{1}^T s = 0$
- node incidence matrix  $A \in \mathbf{R}^{p \times n}$  is

$$A_{ij} = \begin{cases} 1 & \text{arc } j \text{ enters } i \\ -1 & \text{arc } j \text{ leaves node } i \\ 0 & \text{otherwise} \end{cases}$$

- flow conservation:  $Ax + s = 0$
- $\phi(x) = \sum_{j=1}^n \phi_j(x_j)$  is separable convex flow cost function

# Network flow problem

optimal single commodity network flow problem:

$$\begin{array}{ll} \text{minimize} & \sum_{j=1}^n \phi_j(x_j) \\ \text{subject to} & Ax + s = 0 \end{array}$$

- convex, readily solved with standard methods
- dual decomposition yields decentralized solution method

## Network flow Lagrangian

Lagrangian is

$$\begin{aligned} L(x, \nu) &= \phi(x) + \nu^T (Ax + s) \\ &= \nu^T s + \sum_{j=1}^n (\phi_j(x_j) + (a_j^T \nu) x_j) \end{aligned}$$

- $a_j$  is  $j$ th column of  $A$
- we'll interpret  $\nu_i$  as potential at node  $i$
- we use  $\Delta \nu_j$  to denote  $a_j^T \nu$ , which is potential difference across edge  $j$

## Network flow dual

dual function:

$$\begin{aligned} g(\nu) &= \inf_x L(x, \nu) \\ &= \nu^T s + \sum_{j=1}^n \inf_{x_j} (\phi_j(x_j) + (\Delta \nu_j) x_j) \\ &= \nu^T s - \sum_{j=1}^n \phi_j^*(-\Delta \nu_j) \end{aligned}$$

dual problem: maximize  $g(\nu)$



## Recovering primal from dual

- strictly convex  $\phi_j$  means unique minimizer  $x_j^*(y)$  of  $\phi_j(x_j) - yx_j$
- if  $\phi_j$  is differentiable,  $x_j^*(y) = (\phi_j')^{-1}(y)$  (inverse of derivative function)
- optimal flows, from optimal potentials:  $x_j^* = x_j^*(-\Delta\nu_j^*)$
- subgradient of negative dual function:

$$-(Ax^*(\Delta\nu) + s) \in \partial(-g)(\nu)$$

(negative of flow conservation residual)

# Dual decomposition network flow algorithm

**given** initial potential vector  $\nu$ .

**repeat**

1. Determine link flows from potential differences.

$$x_j := x_j^*(-\Delta\nu_j), \quad j = 1, \dots, n.$$

2. Compute flow surplus at each node.

$$S_i := a_i^T x + s_i, \quad i = 1, \dots, p.$$

3. Update node potentials.

$$\nu_i := \nu_i + \alpha_k S_i, \quad i = 1, \dots, p.$$

$\alpha_k$  is an appropriate step size

# Dual decomposition network flow algorithm

- decentralized:
  - flow calculated from potential difference across edge
  - node potential updated from its own flow surplus
- $g(\nu)$  gives lower bound on  $p^*$
- flow conservation  $Ax + s = 0$  only holds in limit

## Electrical network analogy

- electrical network with node incidence matrix  $A$ , nonlinear resistors in branches
- variable  $x_j$  is the current flow in branch  $j$
- source  $s_i$  is external current injected at node  $i$  (must sum to zero)
- flow conservation equation  $Ax + s = 0$  is Kirkhoff Current Law (KCL)
- dual variables are node potentials;  $\Delta\nu_j$  is  $j$ th branch voltage
- branch current-voltage characteristic is  $x_j = x_j^*(-\Delta\nu_j)$

then, current and potentials in circuit are optimal flows and dual variables

## Example: Minimum queueing delay

flow cost function

$$\phi_j(x_j) = \frac{x_j}{c_j - x_j}, \quad \text{dom } \phi_j = [0, c_j)$$

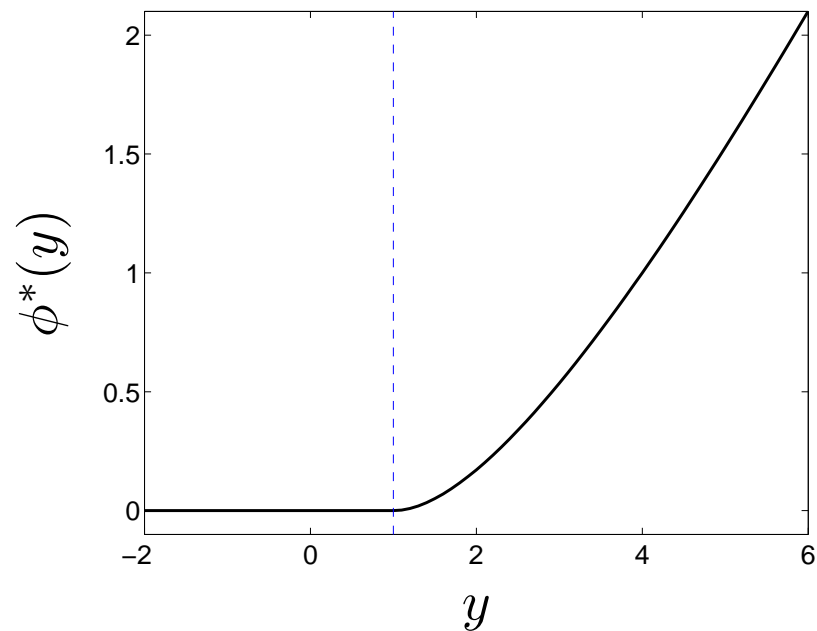
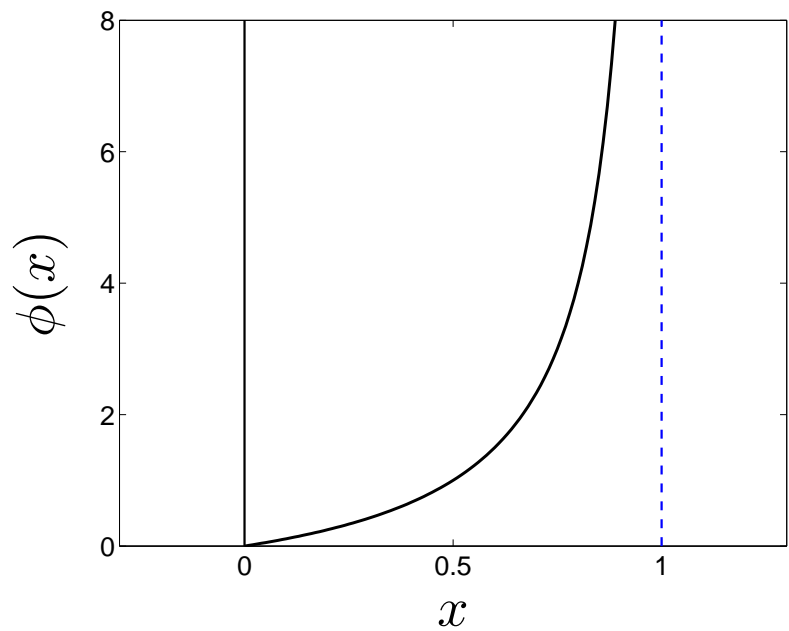
where  $c_j > 0$  are given *link capacities*

( $\phi_j(x_j)$  gives expected waiting time in queue with exponential arrivals at rate  $x_j$ , exponential service at rate  $c_j$ )

conjugate is

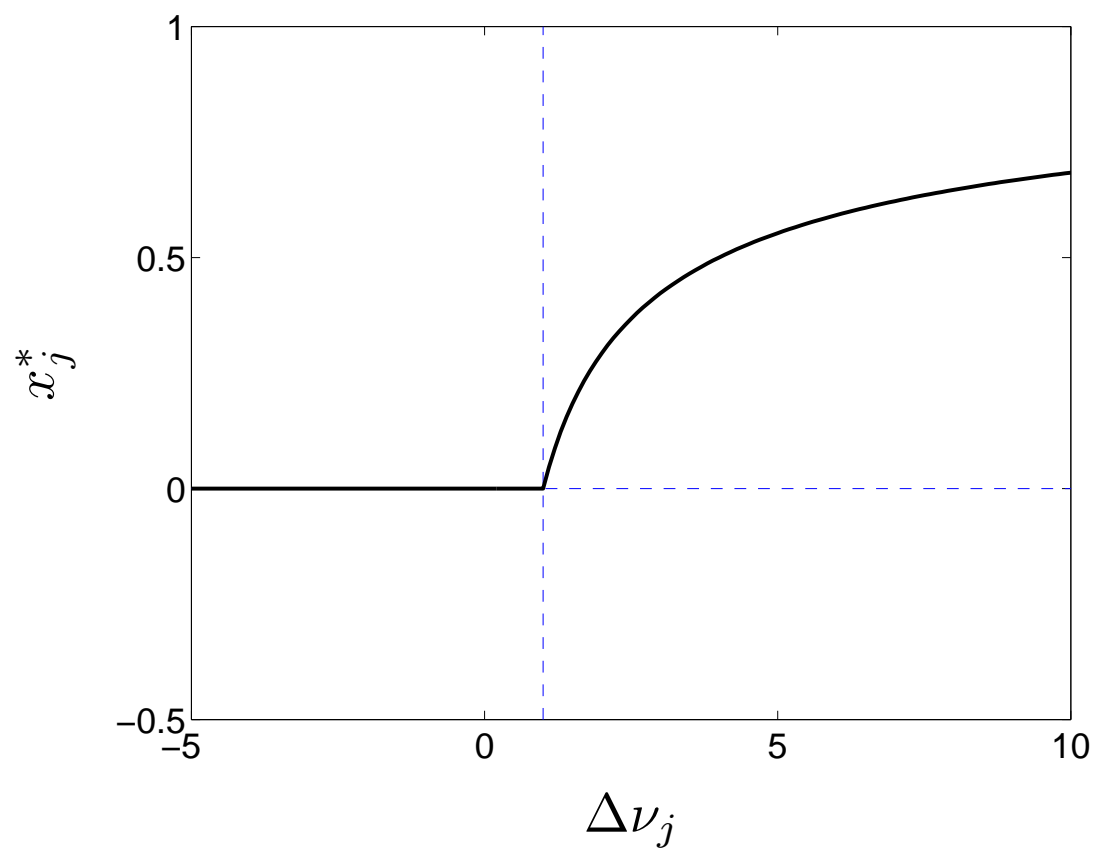
$$\phi_j^*(y) = \begin{cases} (\sqrt{c_j y} - 1)^2 & y > 1/c_j \\ 0 & y \leq 1/c_j \end{cases}$$

cost function  $\phi(x)$  (left) and its conjugate  $\phi^*(y)$  (right),  $c = 1$



(note that conjugate is differentiable)

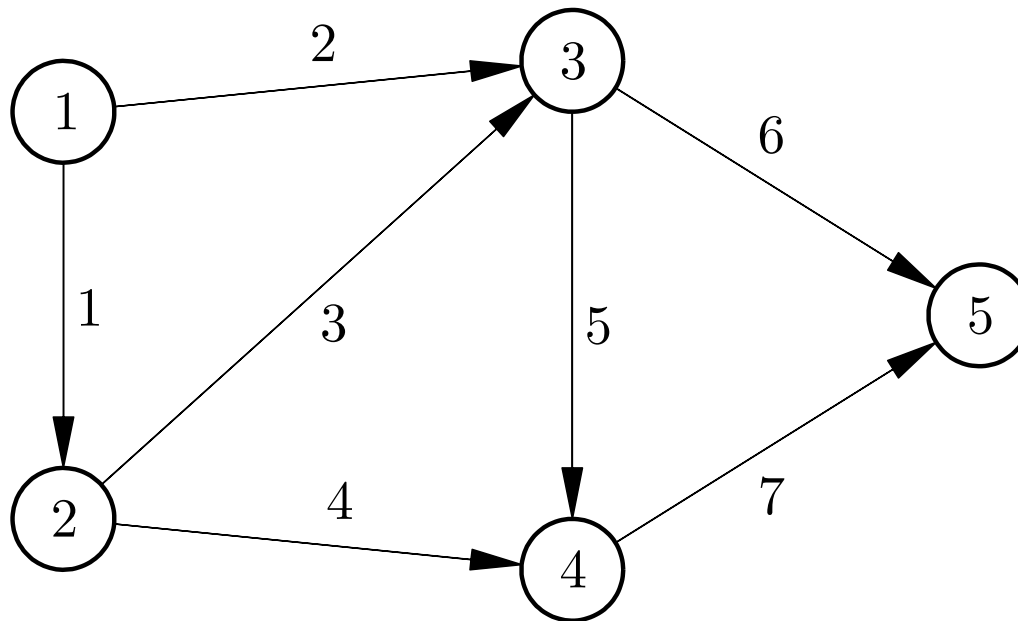
$x_j^*(-\Delta\nu_j)$ , for  $c_j = 1$



gives flow as function of potential difference across link

## A specific example

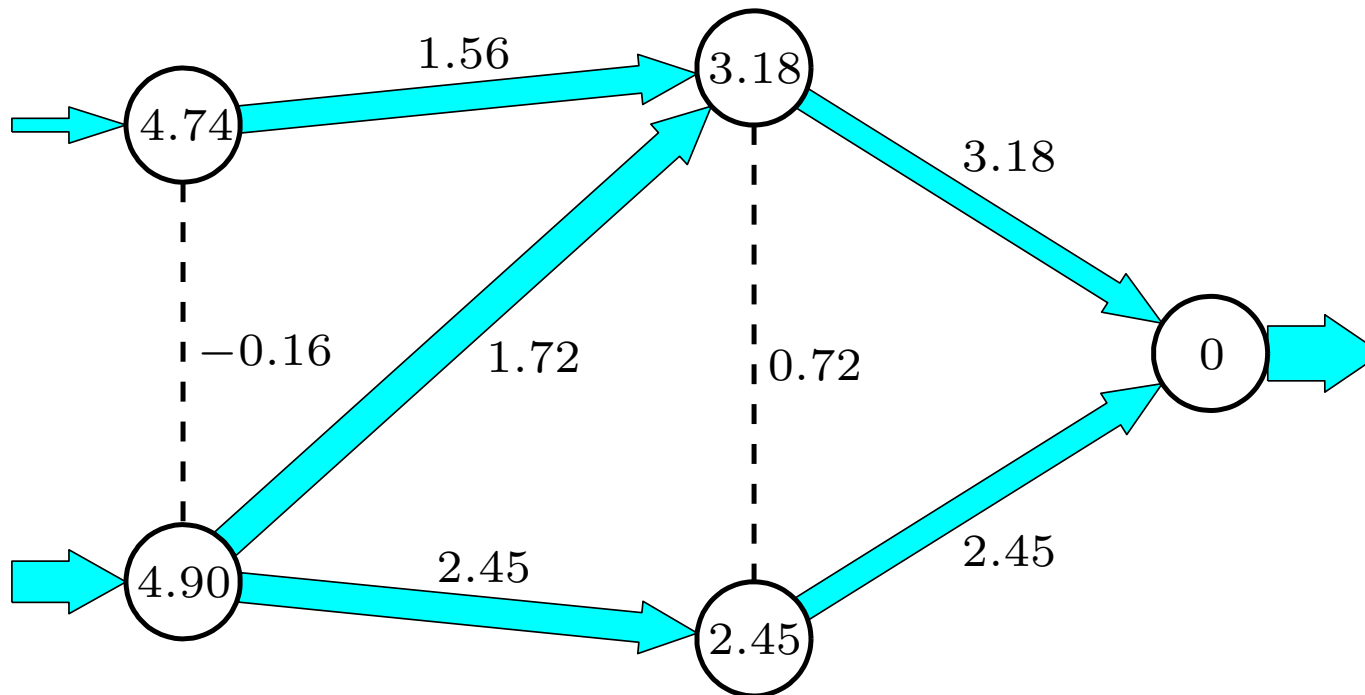
network with 5 nodes, 7 links, capacities  $c_j = 1$





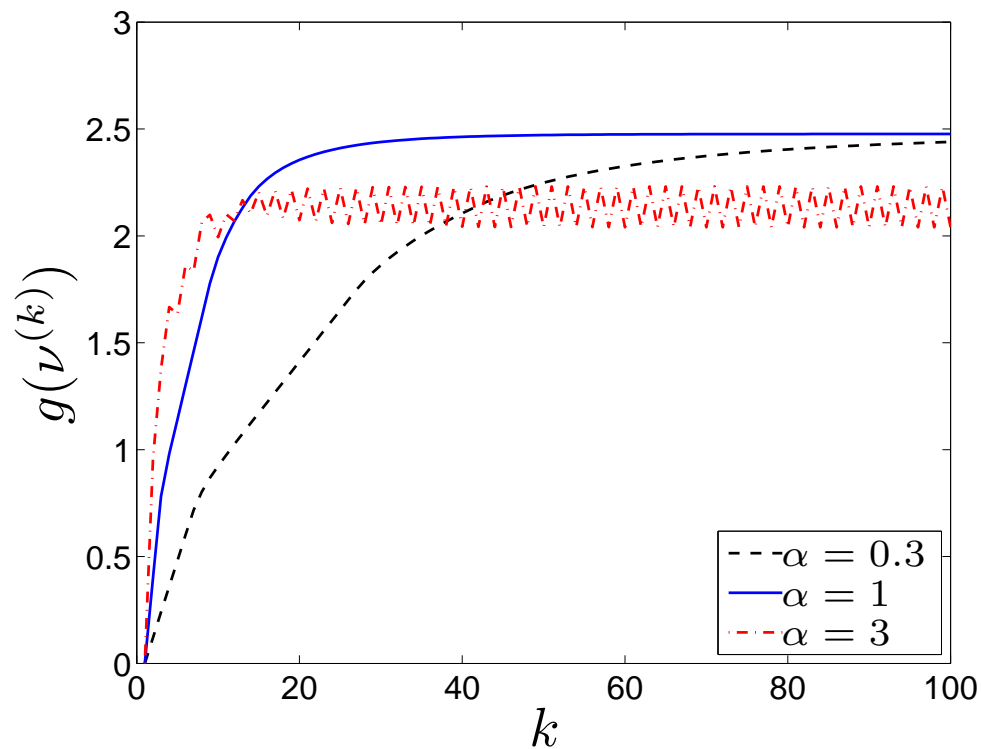
# Optimal flow

optimal flows shown as width of arrows; optimal dual variables shown in nodes; potential differences shown on links



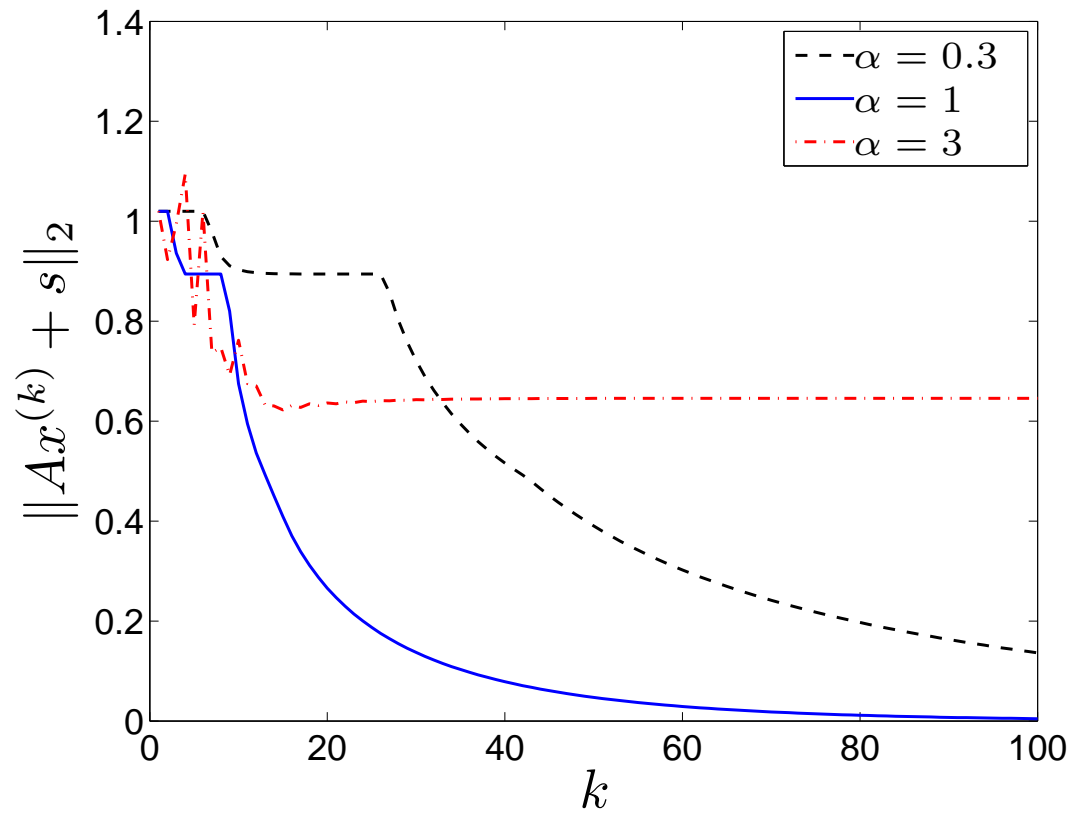
## Convergence of dual function

fixed step size rules,  $\alpha = 0.3, 1, 3$



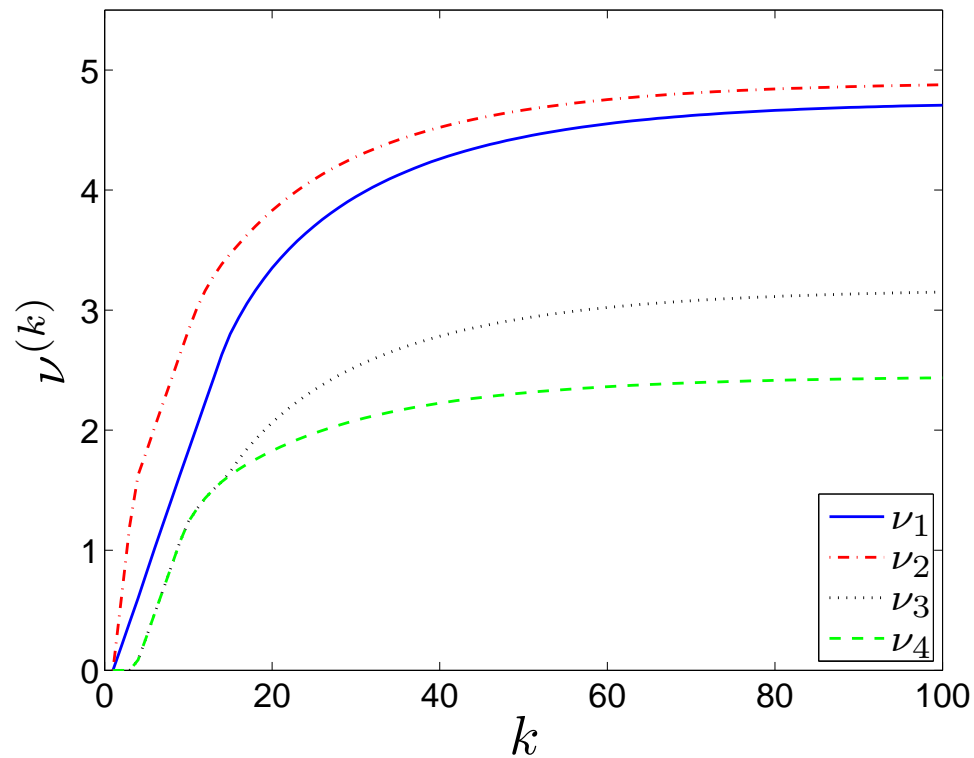
for  $\alpha = 1$ , converges to  $p^* = 2.48$  in around 40 iterations

# Convergence of primal residual



## Convergence of dual variables

$\nu^{(k)}$  versus iteration number  $k$ , fixed step size rule  $\alpha = 1$



( $\nu_5$  is fixed as zero)